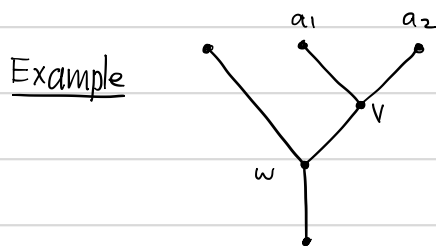


Trees and signs for A_∞ -categories (checked)

We revisit some of the sign issues from ainfmf14 which have come up again in ainfcat3. Throughout S is a commutative, associative but not necessarily unital ring. We begin with basic material on trees.

Definitions Here a tree is a connected acyclic (unoriented) graph. A rooted tree is a tree in which one vertex has been designated the root (for us, the root is always a leaf vertex, i.e. a vertex of valency 1). We make a rooted tree an oriented graph by orienting all edges towards the root (i.e. in the direction of any path to the root). In a rooted tree the parent of a vertex is the vertex connected to it on the path to the root. Every vertex but the root has a parent. A child of a vertex v is a vertex of which v is the parent. A plane tree is a rooted tree together with an ordering for the children of each vertex. A morphism of plane trees is a morphism of oriented graphs which preserves the root vertex and the ordering on children, i.e. if $w \leq w'$ are children of v then $f(w) \leq f(w')$ as children of $f(v)$. A plane tree is valid if it has $n+1$ leaves (including the root) for some $n \geq 2$ and all non-leaves have valency at least three. We call leaves external vertices, non-leaves internal vertices, the edges meeting a leaf are external edges, the others are internal edges.

Defⁿ Given a vertex v in a rooted tree T the depth of v is the length of the unique path from v to the root. The height of v is the length of the longest path from v to a leaf. The height of the tree is the height of the root.



T has height 3, v has depth 2 and height 1, while w has depth 1 and height 2.

Note that our embeddings in \mathbb{R}^2 (= the page) assign the clockwise order (i.e. the children of v are, in order, (a_1, a_2)).

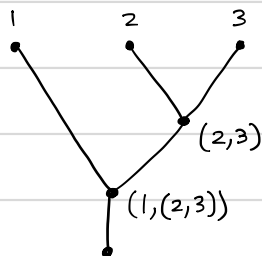
Defⁿ Let T be a plane tree and define a relation $<$ on the vertices as follows:

$w < w'$ if there is a vertex v in $P_w \cap P_{w'}$ (here $P_w, P_{w'}$ are respectively the paths from w, w' to the root) and children c, c' of v with $c < c'$ according to the chosen order, and $c \in P_w, c' \in P_{w'}$.

Lemma If we restrict $<$ to (a) all vertices at a fixed depth or (b) non-root leaves then it is a total order.

Defⁿ Let T be a valid plane tree with $n+1$ leaves. The combinatorial tree of T , denoted $c(T)$, is an ordered set defined as follows. Let the non-root leaves be v_1, \dots, v_k (in order, according to $<$ above) and define $\alpha(v_i) = i$. We define α recursively on non-leaf vertices by the condition that if v has ordered children w_1, \dots, w_n then $\alpha(v) = (\alpha(w_1), \dots, \alpha(w_n))$.

Example



Clearly the set of isomorphism classes of valid plane trees with $n+1$ leaves is in bijection with a set of (suitably defined) combinatorial trees. This set we denote by \mathcal{T}_n .

Example $\mathcal{T}_2 = \left\{ \begin{array}{c} 1 \quad 2 \\ \diagdown \quad \diagup \\ \bullet \end{array} \right\}, \quad \mathcal{T}_3 = \left\{ \begin{array}{c} 1 \quad 2 \quad 3 \\ \diagdown \quad \diagup \quad \diagup \\ \bullet \end{array}, \begin{array}{c} 1 \quad 2 \quad 3 \\ \diagdown \quad \diagdown \quad \diagup \\ \bullet \end{array}, \begin{array}{c} 1 \quad 2 \quad 3 \\ \diagdown \quad \diagdown \quad \diagdown \\ \bullet \end{array} \right\}$

Note that the two trees in \mathcal{T}_3 are isomorphic as oriented graphs, but not as plane trees.

Defⁿ Let S be a commutative, associative but not-necessarily unital ring, and Q a plane tree. A decoration of Q is the following data, where "graded" means either \mathbb{Z} or \mathbb{Z}_2 -graded: (we assume Q has at least one edge)

- (i) a graded S -module L_v for each leaf v (inc. the root),
- (ii) a graded S -module M_e for every edge e ,
- (iii) for each internal vertex v with incoming edges (in order) e_1, \dots, e_k and outgoing edge e an integer N_v (in \mathbb{Z} or \mathbb{Z}_2) and a degree N_v S -linear map

$$\phi_v : M_{e_1} \otimes_S \dots \otimes_S M_{e_k} \longrightarrow M_e. \quad (3.1)$$

- (iv) for each non-root leaf vertex ℓ a degree zero S -linear map

$$\phi_\ell : L_\ell \longrightarrow M_e$$

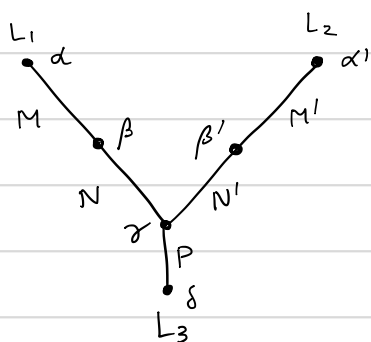
where e is incident at ℓ .

- (v) a degree zero S -linear map

$$\phi_r : M_e \longrightarrow L_r$$

where e is incident at r , which is the root.

Example



$$\begin{aligned} \alpha : L_1 &\rightarrow M, \alpha' : L_2 \rightarrow M', \\ \beta : M &\rightarrow N, \beta' : M' \rightarrow N', \\ \gamma : N \otimes N' &\rightarrow P, \\ \delta : P &\rightarrow L_3 \end{aligned} \quad (3.2)$$

Defⁿ Suppose in \mathcal{Q} that every non-root leaf has equal depth. The height denotation $\langle D \rangle_h$ of D is a homogeneous \mathcal{S} -linear map $L_{e_1} \otimes_{\mathcal{S}} \cdots \otimes_{\mathcal{S}} L_{e_n} \longrightarrow L_r$ where e_1, \dots, e_n denote the non-root leaves (in order), defined as follows:

- Let h be the height of \mathcal{Q} (recall every non-root leaf has depth h). For $1 \leq d \leq h$ let Q_d be the set of vertices of depth d (so Q_h is the set of non-root leaves). Make Q_d a totally ordered set via (6.1), say $Q_d = (v_1, \dots, v_m)$ and for $1 \leq i \leq m$ let $(e_i^1, \dots, e_i^{a_i})$ be the incoming edges at v_i , in their assigned order, and f_i the outgoing edge at v_i . Define a homogeneous \mathcal{S} -linear map for $d < h$ by

$$\Phi_d := \phi_{v_1} \otimes \cdots \otimes \phi_{v_m} : \left(\bigotimes_j M_{e_j^1} \right) \otimes_{\mathcal{S}} \cdots \otimes_{\mathcal{S}} \left(\bigotimes_j M_{e_j^{a_j}} \right) \downarrow M_{f_1} \otimes_{\mathcal{S}} \cdots \otimes_{\mathcal{S}} M_{f_m}. \quad (4.1)$$

using the Koszul sign convention (see p. ⑥ ainfcat2).

- Define Φ_h separately to be $\Phi_h = \bigotimes_{v \in Q_h} \phi_v$ and $\Phi_0 = \phi_r$ where r is the root. Finally,

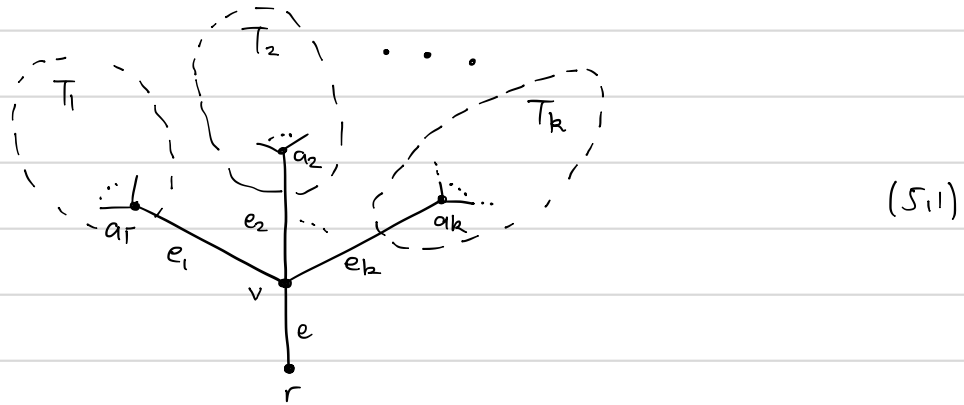
$$\langle D \rangle_h := \Phi_0 \circ \Phi_1 \circ \cdots \circ \Phi_h \quad (4.2)$$

Example For the decoration of (3.2)

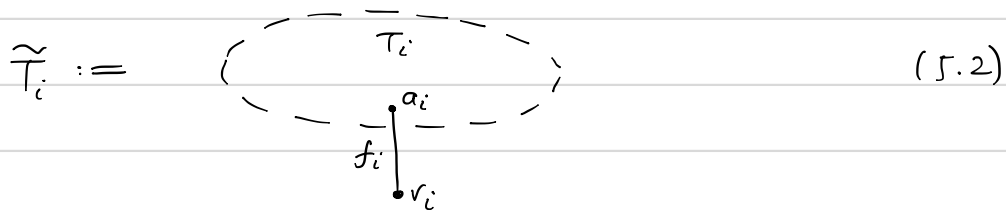
$$\langle D \rangle_h = \underset{\uparrow \Phi_0}{\delta} \circ \underset{\uparrow \Phi_1}{\gamma} \circ \underset{\uparrow \Phi_2}{(\beta \otimes \beta')} \circ \underset{\uparrow \Phi_3}{(\alpha \otimes \alpha')} \quad (4.3)$$

Note that there appears to be a typo in $[L]$ first display after Figure 1, which should read $\text{id}_B \otimes r_2 \otimes r_3 \otimes \text{id}_B$.

Defⁿ Let \mathcal{Q} be any plane tree and D a decoration. Let r be the root and v the vertex x adjacent to r . Suppose v has valence $k+1$ (possibly $k=0$). Consider the diagram (as usual we assume \mathcal{Q} has at least one edge)

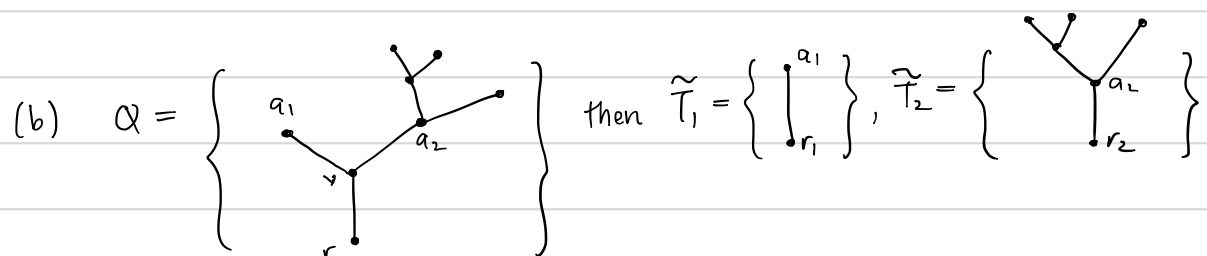


We define plane trees $\tilde{T}_1, \dots, \tilde{T}_k$ to be (note e_i is not in \tilde{T}_i)



where r_i is some new vertex, which we declare the root of \tilde{T}_i . We make \tilde{T}_i a plane tree using the ordering from \mathcal{Q} , and decorate \tilde{T}_i according to a decoration D_i which agrees with D plus assigns Me_i from D to f_i and also $L_{r_i} = Me_i$, $\phi_{r_i} = \text{id}$, ϕ_{a_i} as in D .

Example (a) $\mathcal{Q} = \left\{ \begin{array}{c} v \\ | \\ e \\ | \\ r \end{array} \right\}$ so there are no \tilde{T}_i .



Defⁿ Let Q, D be as above (Q any plane tree with ≥ 1 edge). The branch denotation $\langle D \rangle_B$ is a homogeneous S -linear map $L_{e_1} \otimes_S \cdots \otimes_S L_{e_n} \longrightarrow L_r$ where e_1, \dots, e_n denote the non-root leaves (in order), defined as follows:

- if Q has one edge, so

$$Q = \left\{ \begin{array}{c} L \\ \downarrow e \\ v \\ \downarrow e \\ r \\ L_r \end{array} \right\}, \quad \phi_v: L \rightarrow Me, \quad \phi_r: Me \rightarrow L_r \quad (6.1)$$

then $\langle D \rangle_B := \phi_r \circ \phi_v$.

- otherwise we define recursively using (5.1)

$$\langle D \rangle_B := \phi_r \circ \phi_v \circ (\langle D_1 \rangle_B \otimes_S \cdots \otimes_S \langle D_k \rangle_B). \quad (6.2)$$

where D_i is the decoration of \tilde{T}_i defined above.

Remark We may write $(e_1, \dots, e_n) = (e_1^{m_1}, \dots, e_1^{m_{n_1}}, e_2^{m_2}, \dots, e_2^{m_{n_2}}, \dots)$ where e_i^j are the non-root leaves in T_i , so that $\langle D_i \rangle_B: L_{e_i^1} \otimes_S \cdots \otimes_S L_{e_i^{m_i}} \longrightarrow Me_i$. Hence $\langle D_1 \rangle_B \otimes \cdots \otimes \langle D_k \rangle_B: L_{e_1} \otimes \cdots \otimes L_{e_n} \longrightarrow Me_1 \otimes \cdots \otimes Me_k$. Again, the Koszul sign rule is used to define (6.2).

Defⁿ Let Q be a plane tree, D a decoration, and L_1, \dots, L_n, L_r the modules assigned to the leaves so that $\langle D \rangle_B: L_1 \otimes_S \cdots \otimes_S L_n \longrightarrow L_r$. We define

$$\text{eval}_D: L_1 \otimes_S \cdots \otimes_S L_n \longrightarrow L_r$$

to be the S -linear map associated to D as a diagram in $S\text{Mod}$, ignoring the grading. Recursively, if Q has one edge $\text{eval}_D = \phi_r \circ \phi_v$ and in the case of (6.2) we use the same formula but as plain S -linear maps (i.e. no Koszul signs).

Lemma With the above notation

$$\langle D \rangle_B(a_1 \otimes \dots \otimes a_n) = (-1)^s \text{eval}_D(a_1 \otimes \dots \otimes a_n) \quad (6.5.1)$$

where we write ℓ_1, \dots, ℓ_n for the non-root leaves in their order, and

$$s = \sum_{i=1}^n \sum_{v \succ \ell_i} |a_i| \cdot |\phi_v| \quad (6.5.2)$$

where v ranges over all vertices, including leaves.

Proof By induction on the height of Q . If the height is 1 we are in the situation of (6.1), and the claim is clear. Otherwise

$$\begin{aligned} \langle D \rangle_B(a_1 \otimes \dots \otimes a_n) &= \phi_r \circ \phi_v \circ (\langle D_1 \rangle_B \otimes_s \dots \otimes_s \langle D_k \rangle_B)(a_1 \otimes \dots \otimes a_n) \\ &= (-1)^t \phi_r \circ \phi_v (\langle D_1 \rangle_B(a_1 \otimes \dots \otimes a_{m_1}) \\ &\quad \otimes \langle D_2 \rangle_B(a_{m_1+1} \otimes \dots \otimes a_{m_2}) \\ &\quad \otimes \dots \\ &\quad \otimes \langle D_k \rangle_B(a_{m_{k-1}+1} \otimes \dots \otimes a_n)) \end{aligned}$$

where (set $m_0 = 0$)

$$t = \sum_{i=1}^k \left(\sum_{j=m_{i-1}+1}^{m_i} |a_j| \right) \left(\sum_{\ell=i+1}^k |\langle D_\ell \rangle_B| \right)$$

By the inductive hypothesis each $\langle D_i \rangle_B(\dots)$ can be written as a sign times $\text{eval}_{D_i}(\dots)$, and is easy to see the signs match. \square

Proof By induction on h , the height of Q . The base case is $h=1$, i.e.

$$Q = \left\{ \begin{array}{c} v \\ \vdots \\ e \\ \vdots \\ r \end{array} \right\}$$

in which case the claim is trivially true. Now suppose $h > 1$ and that we decompose Q as in (5.1). Then by def^N of $\langle D \rangle_B$ and the inductive hypothesis

$$\begin{aligned} \langle D \rangle_B &= \phi_r \circ \phi_v \circ (\langle \tilde{T}_1 \rangle_B \otimes_s \cdots \otimes_s \langle \tilde{T}_k \rangle_B) \\ &= \phi_r \circ \phi_v \circ ((-1)^{\mathcal{T}(\tilde{T}_1)} \langle \tilde{T}_1 \rangle_H \otimes \cdots \otimes (-1)^{\mathcal{T}(\tilde{T}_k)} \langle \tilde{T}_k \rangle_H) \quad (8.2) \\ &= (-1)^{\sum_i \mathcal{T}(\tilde{T}_i)} \phi_r \circ \phi_v \circ (\langle \tilde{T}_1 \rangle_H \otimes \cdots \otimes \langle \tilde{T}_k \rangle_H) \end{aligned}$$

where $\mathcal{T}(\tilde{T}_i)$ denotes the sign (7.2) for the induced decoration of \tilde{T}_i . Note that by our hypothesis on Q , all the \tilde{T}_i have height $h-1$, and so by def^N

$$\langle \tilde{T}_i \rangle_H = \Phi_2^i \circ \Phi_3^i \circ \cdots \circ \Phi_h^i \quad \text{where} \quad \Phi_d^i = \bigotimes_v \phi_v$$

for v ranging over vertices in \tilde{T}_i of depth d in Q (note that $\Phi_2^i = \phi_{a_i}$ in Q).

Observe that $|\Phi_d^i| = \sum_v |\phi_v|$ where v ranges over the contributing vertices in \tilde{T}_i .

Hence we calculate

$$\begin{aligned} \langle \tilde{T}_1 \rangle_H \otimes \cdots \otimes \langle \tilde{T}_k \rangle_H &= (\Phi_2^1 \circ \cdots \circ \Phi_h^1) \otimes (\Phi_2^2 \circ \cdots \circ \Phi_h^2) \otimes \cdots \\ &\quad \cdots \otimes (\Phi_2^k \circ \cdots \circ \Phi_{h-1}^k) \\ &= (-1)^s (\Phi_2^1 \otimes \cdots \otimes \Phi_2^k) \circ (\Phi_3^1 \otimes \cdots \otimes \Phi_3^k) \circ \cdots \quad (8.3) \\ &\quad \cdots \circ (\Phi_h^1 \otimes \cdots \otimes \Phi_h^k) \end{aligned}$$

where $s = \sum_{a < b, i < j} |\Phi_j^a| |\Phi_i^b|$ where $1 \leq a < b \leq k$ and $2 \leq i < j \leq h$

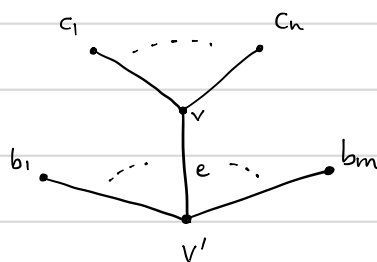
It follows that, combining (8.2) and (8.3),

$$\langle D \rangle_B = \sum_i (-1)^{\mathcal{T}(\tilde{T}_i) + s} \langle D \rangle_H. \quad (9.1)$$

Given a pair of vertices $w < w'$ in Q with $\text{depth}(w') < \text{depth}(w)$ we have $w \in T_i$ and $w' \in T_j$ for some i, j . In fact we must have $i \leq j$. If $i = j$ then $|\phi_w| |\phi_{w'}|$ contributes to $\mathcal{T}(\tilde{T}_i)$ and if $i < j$ then $|\phi_w| |\phi_{w'}|$ contributes to s via $|\Phi_{\text{depth}(w)}^i| |\Phi_{\text{depth}(w')}^j|$, which proves that the sign in (9.1) agrees with (7.2). \square

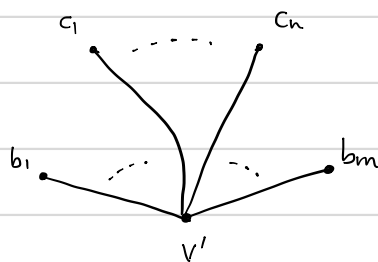
Edge contraction (write $V(T)$ for the vertex set)

Defⁿ Let T be a plane tree and e an internal edge of T , which connects vertices v and v' . Suppose v' is closer to the root and that v has children c_1, \dots, c_n (note $n \geq 1$ since e is internal) while v' has children b_1, \dots, b_m with $b_j = v$.



(10.1)

We define the plane tree $T \setminus e$ to have vertices $V(T) \setminus \{v\}$. We connect c_1, \dots, c_n to v' and order the children of v' as follows: $b_1, \dots, b_{j-1}, c_1, \dots, c_n, b_{j+1}, \dots, b_m$.



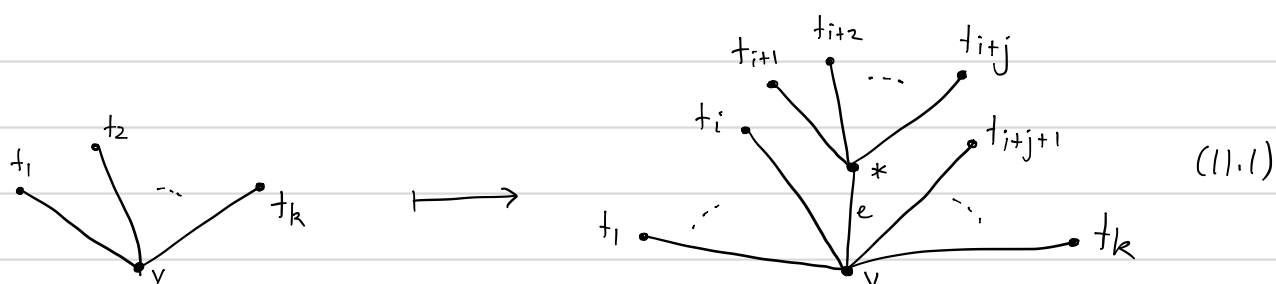
(10.2)

Clearly $T \setminus e$ is a plane tree with the same root as T and one less internal edge.

Edge insertion

Defⁿ Let T be a plane tree and v an internal vertex, with children t_1, \dots, t_k .

Given integers $i \geq 0$ and $j \geq 1$ with $i+j \leq k$ we define the edge insertion $\text{ins}(T, v, i, j)$ to be the plane tree with vertices $V(T) \cup \{*\}$. The parent/child relations (which clearly determine the edges) between vertices in $V(T)$ are preserved in $\text{ins}(T, v, i, j)$ with the exception of t_{i+1}, \dots, t_{i+j} which now have parent $*$. The parent of $*$ is v , and it has t_{i+1}, \dots, t_{i+j} as its only children. Denote the $* - v$ edge by e .



The ordering on children is as in T , with the exception of v and $*$ where (11.1) gives the ordering. Clearly $\text{ins}(T, v, i, j)$ is a plane tree with one more internal edge than T .

Lemma If T is a valid plane tree and e an internal edge then $T \dashv e$ is valid.

If Q is a valid plane tree, v an internal vertex with k children, and $2 \leq j \leq k-1$ then $\text{ins}(Q, v, i, j)$ is valid.

Proof Note that in $\text{ins}(Q, v, i, j)$ the vertex v has $k-j+1$ children, so for the tree to be valid we need both $j \geq 2$ and $j \leq k-1$. \square

Let \mathcal{C}_n be the category of valid plane trees with $n+1$ leaves, so that \mathcal{T}_n is the set of isomorphism classes of objects of \mathcal{C}_n . We can write $\mathcal{T}_n = \coprod_{c \geq 0} \mathcal{T}_n^{(c)}$ where $\mathcal{T}_n^{(c)}$ denotes trees with c internal edges, i.e. $e_i(T) = c$.

Defⁿ \mathcal{T}_n^+ is the set of pairs (T, e) where $T \in \mathcal{T}_n$ and e is an internal edge of T . (note that $\mathcal{T}_2^+ = \emptyset$ since $\mathcal{T}_2 = \{\text{Y}\}$).

Let us write $|v|$ for the valency of v (so if v has k children, $|v| = k+1$).

FALSE

Lemma For $n > 2$ the map $\mathcal{T}_n^+ \xrightarrow{\pi} \mathcal{T}_n$, $\pi(T, e) = T \dashv e$ is surjective, and

$$\pi^{-1}(Q) = \{ (\text{ins}(Q, v, i, j), e) \mid v \text{ an internal vertex, } i \geq 0, i+j \leq |v|-1, 2 \leq j \leq |v|-2 \}$$

Proof If $Q \in \mathcal{T}_n$ then since $n \geq 2$ there is at least one internal vertex. Clearly $\text{ins}(Q, v, i, j) \sqsupset e = Q$ so π is surjective. Suppose that $T \sqsupset e = Q$. In the situation of (10.1), (10.2) we have

$$T = \text{ins}(Q, v', j-1, n)$$

proving the second statement. \square

This implies there is a bijection between pairs $(T, e) \in \mathcal{T}_n^+$ and tuples (Q, v, i, j) with $Q \in \mathcal{T}_n$, v an internal vertex, $i \geq 0$, $i+j \leq |v|-1$, $2 \leq j \leq |v|-2$ for any $n \geq 2$. More precisely: (noting $|v| = k+1$)

Lemma For $n \geq 2$ and $c \geq 0$ there is a bijection

$$(\mathcal{T}_n^{(c+1)})^+ \longrightarrow \left\{ (Q, v, i, j) \mid Q \in \mathcal{T}_n^{(c)}, v \text{ an internal vertex}, \right. \\ \left. i \geq 0, i+j \leq |v|-1, 2 \leq j \leq |v|-2 \right\}. \quad (12.1)$$

Note We restrict to $n \geq 2$ because $\mathcal{T}_2^+ = \emptyset$.