

Monoidal Categories:

Important Example:

Consider Vec _{\mathbb{R}}

Then the tensor product

$$\otimes_{\mathbb{R}} : \text{Vec}_{\mathbb{R}} \times \text{Vec}_{\mathbb{R}} \rightarrow \text{Vec}_{\mathbb{R}}$$

and $k \in \text{Vec}_{\mathbb{R}}$ define a monoidal category.

Note: Algebraic monoid is a set with an associative product and a unit.

Def^m: (Monoidal Category)

A (strict) monoidal category is a category \mathcal{V} with two functors

Outline:

- ① Monoidal Cat
- ② Monoids in Monoidal Cat
- ③ Δ^+ & Monoidal Functors
- ④ Symmetric Monoidal Cat
- ⑤ Symmetric Monoidal Functors & \mathbb{I}
- ⑥ Frobenius Objects.

$$\mu: V \times V \rightarrow V$$

(eg $\otimes_{\mathbb{K}}: \text{Vec}_{\mathbb{K}} \times \text{Vec}_{\mathbb{K}} \rightarrow \text{Vec}_{\mathbb{K}}$)

$$\eta: \mathbb{1} \rightarrow V$$

(eg: $\eta(\mathbb{1}) = \mathbb{K}$)

category with one object and one morphism

such that the following diagrams commute

$$\begin{array}{ccc} V \times V \times V & & \\ \swarrow \mu \times \text{id}_V & & \searrow \text{id}_V \times \mu \\ V \times V & & V \times V \\ \searrow \mu & & \swarrow \mu \\ V & & V \end{array}$$

eg (A, B, C)

$$\begin{array}{ccc} & & \\ \swarrow & & \searrow \\ (A \otimes_{\mathbb{K}} B, C) & & (A, B \otimes_{\mathbb{K}} C) \\ \searrow & & \swarrow \\ (A \otimes_{\mathbb{K}} B) \otimes_{\mathbb{K}} C \cong A \otimes_{\mathbb{K}} (B \otimes_{\mathbb{K}} C) & & \end{array}$$

$$\begin{array}{ccc} \mathbb{1} \times V & \xrightarrow{\eta \times \text{id}_V} & V \times V \\ & \searrow & \downarrow \mu \\ & & V \end{array}$$

$$\begin{array}{ccc} (\mathbb{1}, A) & \rightarrow & (\mathbb{K}, A) \\ & \searrow & \downarrow \\ & & A \cong \mathbb{K} \otimes_{\mathbb{K}} A \end{array}$$

For strict monoidal categories these are equalities

$$\begin{array}{ccc} V \times V & \xleftarrow{\text{id}_V \times \eta} & V \times \mathbb{1} \\ \downarrow \mu & & \swarrow \\ V & & \end{array}$$

$$\begin{array}{ccc} (A, \mathbb{K}) & \leftarrow & (A, \mathbb{1}) \\ \downarrow & & \swarrow \\ A \otimes_{\mathbb{K}} \mathbb{K} \cong A & & \end{array}$$

Functoriality of μ :

For a monoid category (V, μ, η)
and $A, B \in \text{ob}(V)$ we denote

$$\mu(A, B) = A \otimes B \quad (\text{or } A \otimes B, A \sqcup B)$$

Functoriality of μ means that for

$$A \xrightarrow{f} A' \xrightarrow{f'} A'' \quad B \xrightarrow{g} B' \xrightarrow{g'} B''$$

we have

$$A \otimes B \xrightarrow{f \otimes g} A' \otimes B' \xrightarrow{f' \otimes g'} A'' \otimes B''$$

$\xrightarrow{(f' \circ f) \circ (g' \circ g)}$

And $\text{id}_{A \otimes B} = \text{id}_A \otimes \text{id}_B$

Examples: 1) $(\text{Set}, \times, \{*\})$ $\left(\begin{array}{l} (A \times B) \times C \cong A \times (B \times C) \\ \{*\} \times A \cong A \cong A \times \{*\} \end{array} \right)$

2) $(\text{Set}, \sqcup, \emptyset)$ $\left(\begin{array}{l} (A \sqcup B) \sqcup C \cong A \sqcup (B \sqcup C) \\ \emptyset \sqcup A \cong A \cong A \sqcup \emptyset \end{array} \right)$

3) $(\text{Vec}_{\mathbb{k}}, \otimes, \mathbb{k})$

4) $(2\text{Cob}, \parallel, \emptyset)$ — will see in detail next week

$$\text{Ob}(2\text{Cob}) = \text{closed oriented 1-mflds} \\ = \{ \underbrace{S' \parallel \dots \parallel S'}_{n \text{ times}} : n \in \mathbb{Z}_{\geq 0} \}$$

$$\text{Hom}(\underbrace{S' \parallel \dots \parallel S'}_{m \text{ times}}, \underbrace{S' \parallel \dots \parallel S'}_{n \text{ times}}) = \text{closed oriented 2-mflds } M \text{ with} \\ \partial M = \underbrace{(S' \parallel \dots \parallel S')}_{m \text{ times}} \parallel \underbrace{(S' \parallel \dots \parallel S')}_{n \text{ times}}$$

Note: there is a sense of "in" and "out" boundaries which formalized by fixing embeddings $f_{\text{in}}: (S' \parallel \dots \parallel S')^* \rightarrow \partial M$ and $f_{\text{out}}: (S' \parallel \dots \parallel S') \rightarrow \partial M$. reversed orientation

eg of composition

$$S', S' \parallel S', \emptyset \in \text{Ob}(2\text{Cob}) \text{ ie 1-mflds}$$

$$f: S' \rightarrow S' \parallel S' = \text{cylinder}, \quad g: S' \parallel S' \rightarrow \emptyset = \text{pair of pants}$$

$$g \circ f = \text{cylinder with hole} = \text{annulus}$$

$$\left(\begin{array}{c} \text{S} \\ \text{S} \end{array} \right) \circ \left(\begin{array}{c} \text{S} \\ \text{S} \end{array} \right) = \left(\begin{array}{cc} \text{S} & \text{S} \\ \text{S} & \text{S} \end{array} \right)$$

Here functoriality of \amalg follows from the following pictures.

$$\begin{aligned} \text{If } f: S' \rightarrow S' \amalg S' &= \text{cone with two holes} \\ g: S' \rightarrow S' &= \text{cylinder} \end{aligned}$$

$$\begin{aligned} h: S' \amalg S' \rightarrow \emptyset &= \text{cup} \\ k: S' \rightarrow S' \amalg S' \amalg S' &= \text{cone with two holes and a hole below} \end{aligned}$$

$$\text{Then } h \circ f = \text{cup with two holes} \qquad k \circ g = \text{cylinder with two holes and a hole below}$$

$$\text{So } h \circ f \amalg k \circ g = \text{cup with two holes and cylinder with two holes and a hole below}$$

$$\text{Now } f \amalg g = \text{cone with two holes and cylinder} \qquad h \amalg k = \text{cup with two holes and cone with two holes and a hole below}$$

$$(f \amalg g) \circ (h \amalg k) = \text{cone with two holes and cylinder with two holes and a hole below (crossed out)}$$

4) Categories with $\left\{ \begin{array}{l} \text{products and a terminal object} \\ \text{coproducts and an initial object} \end{array} \right.$

$$5) \Delta^+ = \Delta \cup \{\emptyset\}$$

simplex
category \rightarrow Δ
 \downarrow

$$\text{ob}(\Delta^+) = \{ \langle n \rangle \}_{n \in \mathbb{Z}_{\geq 0}} \quad \text{for } n > 0 \quad \langle n \rangle = [n-1]$$

where $\langle n \rangle = (0, 1, \dots, n-1)$ ordered set.

$\text{Hom}(\langle n \rangle, \langle m \rangle) =$ order preserving maps

Let $\otimes: \Delta^+ \times \Delta^+ \rightarrow \Delta^+$ s.t that

$$\langle n \rangle \otimes \langle m \rangle = \langle n+m \rangle$$

For $f: \langle n_1 \rangle \rightarrow \langle m_1 \rangle$ order preserving
 $g: \langle n_2 \rangle \rightarrow \langle m_2 \rangle$

$f \otimes g: \langle n_1+n_2 \rangle \rightarrow \langle m_1+m_2 \rangle$ s.t

$$f \otimes g(x) = \begin{cases} f(x) & \text{if } 0 \leq x \leq n_1 - 1 \\ m_1 + g(x - n_1) & \text{if } n_1 \leq x \leq n_1 + n_2 \end{cases}$$

Then $(\mathbb{A}^+, \otimes, \langle 0 \rangle)$ is a monoidal category.

proof: The main thing to check is functoriality of \otimes .

$$\text{Let } \begin{array}{l} f: \langle n_1 \rangle \rightarrow \langle n_2 \rangle \\ h: \langle m_1 \rangle \rightarrow \langle m_2 \rangle \end{array} \quad \begin{array}{l} g: \langle n_2 \rangle \rightarrow \langle n_3 \rangle \\ k: \langle m_2 \rangle \rightarrow \langle m_3 \rangle \end{array}$$

$$(g \circ f) \otimes (k \circ h): \langle n_1 \rangle \otimes \langle m_1 \rangle \rightarrow \langle n_3 \rangle \otimes \langle m_3 \rangle$$

$$\text{s.t. } (g \circ f) \otimes (k \circ h)(x) = \begin{cases} (g \circ f)(x) & \text{if } 0 \leq x \leq n_1 - 1 \\ n_3 + (k \circ h)(x - n_1) & \text{if } n_1 \leq x \leq n_1 + m_1 \end{cases}$$

$$f \otimes h(x) = \begin{cases} f(x) & \text{if } 0 \leq x \leq n_1 - 1 \\ n_2 + h(x - n_1) & \text{if } n_1 \leq x \leq n_1 + m_1 \end{cases}, \quad g \otimes k(x) = \begin{cases} g(x) & \text{if } 0 \leq x \leq n_2 - 1 \\ n_3 + k(x - n_2) & \text{if } n_2 \leq x \leq n_2 + m_2 \end{cases}$$

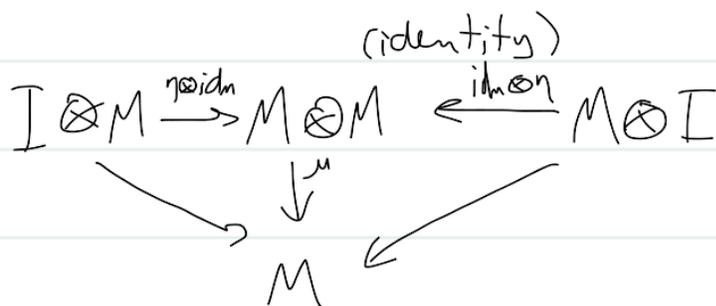
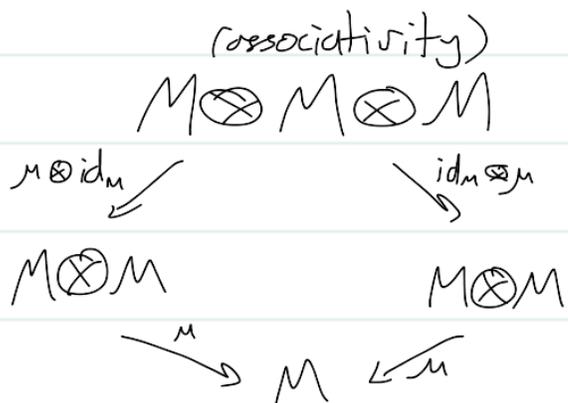
$$(g \otimes k) \circ (f \otimes h)(x) = (g \otimes k) \begin{cases} f(x) & \text{if } 0 \leq x \leq n_1 - 1 \\ n_2 + h(x - n_1) & \text{if } n_1 \leq x \leq n_1 + m_1 \end{cases} = \begin{cases} g \circ f(x) & \text{if } 0 \leq x \leq n_1 - 1 \\ n_3 + g(n_2 + h(x - n_1) - n_2) & \text{if } n_1 \leq x \leq n_1 + m_1 \end{cases}$$

or $\begin{cases} f(x) \leq n_2 - 1 \\ n_2 \leq n_2 + h(x - n_1) \end{cases}$

Def^m: (Monoid (in a monoidal category))

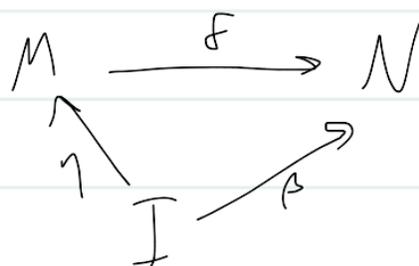
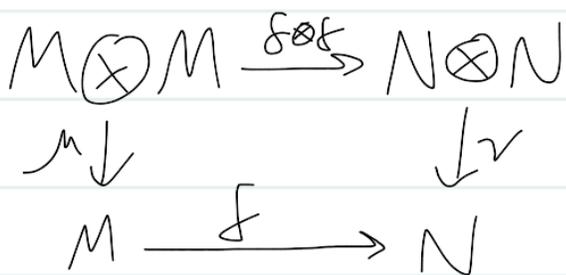
Let (V, \otimes, I) be a monoidal category.

A monoid (M, μ, η) in \mathcal{V} is an object $M \in \text{ob}(\mathcal{V})$ and morphisms $\mu \in \text{Hom}(M \otimes M, M)$ $\eta \in \text{Hom}(I, M)$ s.t the following diagrams commute



Defⁿ (Monoid Morphism)

For monoids $(M, \mu, \eta), (N, \nu, \beta)$ in \mathcal{V} , $f \in \text{Hom}(M, N)$ is a morphism of monoids if



Ex: Mon \subseteq Set
Ring \subseteq Ab

monoids in sets
rings in abelian groups

Monoidal Functors:

For (V, \otimes, I) , (V', \otimes', I') monoidal categories a monoidal functor from V to V' is a functor $F: V \rightarrow V'$ s.t the following diagrams commute.

$$\begin{array}{ccc} V \times V & \xrightarrow{F \times F} & V' \times V' \\ \downarrow \eta & & \downarrow \eta' \\ V & \xrightarrow{F} & V' \end{array}$$

$$\begin{array}{ccc} V & \xrightarrow{F} & V' \\ \eta \swarrow & & \nearrow \eta' \\ & I & \end{array}$$

Example: $F: \text{Vec}_{\mathbb{R}} \rightarrow \text{Vec}_{\mathbb{C}}$
 $A \longmapsto \mathbb{C} \otimes_{\mathbb{R}} A$

Monoidal as $\mathbb{C} \otimes_{\mathbb{R}} (A \otimes_{\mathbb{R}} B) \cong (\mathbb{C} \otimes_{\mathbb{R}} A) \otimes_{\mathbb{C}} (\mathbb{C} \otimes_{\mathbb{R}} B)$

$$\mathbb{C} \otimes_{\mathbb{R}} \mathbb{R} \cong \mathbb{C}$$

Example: Let (V, \otimes, I) be a monoidal category.

Let (M, μ, η) be a monoid in V .

We have a monoidal functor given by

$$F: \mathbb{A}^+ \rightarrow V \quad \text{s.t.} \quad F(\langle n \rangle) = M^{\otimes n}$$

and for $f: \langle n \rangle \rightarrow \langle m \rangle$ we can write f as a combination of $z: \langle 0 \rangle \rightarrow \langle 1 \rangle$, $m: \langle 2 \rangle \rightarrow \langle 1 \rangle$ and $+$.

Therefore defining $F(z) = \eta$ and $F(m) = \mu$ determines $F(f)$ and therefore a functor F .

Likewise given monoidal $F: \mathbb{A}^+ \rightarrow V$ ($F(\langle 1 \rangle), F(z), F(m)$ is a monoid in V).

Def: (Category of Monoids in V)

Let (V, \otimes, I) be a monoidal category.

Let $\text{ob}(\text{Mon}(V))$ be the monoids in V .

For $M, N \in \text{ob}(\text{Mon}(V))$ let $\text{Hom}_{\text{Mon}(V)}(M, N)$ be morphisms of monoids.

Mon(V) is a category.

Theorem 1: $\text{MonCat}(\Delta^+, V) \xleftrightarrow[\text{cat}]{\text{equiv. of}} \text{Mon}(V)$
(monoidal functors from Δ^+ to V)

Remark: Δ^+ is therefore the "free monoidal cat on a single object".

Remark: We have avoided or some technicalities by avoiding the issues around "strictness".

This can all be dealt with but we can ignore these issues for the most part.

See 3.2. 17-19 of Kock for a few details.

Defⁿ: (Symmetric Monoidal Category)

A (strict) monoidal category (V, \otimes, I) with maps $\tilde{\tau}_{x,y}: x \otimes y \rightarrow y \otimes x$ for each $x, y \in \text{ob}(V)$ is called a symmetric monoidal category if

i) the maps $\tilde{\tau}_{x,y}$ are natural i.e.

$$\begin{array}{ccc} X \otimes Y & \xrightarrow{\tilde{\tau}_{x,y}} & Y \otimes X \\ \downarrow f \otimes g & & \downarrow g \otimes f \\ X' \otimes Y' & \xrightarrow{\tilde{\tau}_{x',y'}} & Y' \otimes X' \end{array} \quad \text{commutes}$$

ii) For $x, y, z \in \text{ob}(V)$

$$\begin{array}{ccc} X \otimes Y \otimes Z & \xrightarrow{\tilde{\tau}_{x,y \otimes z}} & Y \otimes Z \otimes X \\ \tilde{\tau}_{x,y} \square \text{id}_z \searrow & \nearrow \text{id}_y \square \tilde{\tau}_{x,z} & \\ Y \otimes X \otimes Z & & \end{array} \quad \begin{array}{ccc} X \otimes Y \otimes Z & \xrightarrow{\tilde{\tau}_{x \otimes y, z}} & Z \otimes X \otimes Y \\ \text{id}_x \square \tilde{\tau}_{y,z} \searrow & \nearrow \tilde{\tau}_{x,z} \square \text{id}_y & \\ X \otimes Z \otimes Y & & \end{array}$$

iii) $\tilde{\tau}_{y,x} \circ \tilde{\tau}_{x,y} = \text{id}_{x \otimes y}$

Example: 1) $A \otimes B \xrightarrow{\tau_{A,B}} B \otimes A$ the linear map such that $\tau_{A,B}(a \otimes b) = b \otimes a$

2) Previous examples such as set with $x, 1$

3) $\text{gr-Vec}_{\mathbb{K}} = \text{graded vector spaces}$

for $V = \bigoplus_{n \in \mathbb{Z}} V_n$ and $W = \bigoplus_{n \in \mathbb{Z}} W_n$

$$(V \otimes W)_n = \bigoplus_{p+q=n} V_p \otimes W_q$$

$$I = \bigoplus_{n \in \mathbb{Z}} I_n \quad \text{with} \quad I_n = \begin{cases} \mathbb{K} & \text{if } n=0 \\ 0 & \text{o.w.} \end{cases}$$

Can define two twist maps for $v \in V_p, w \in W_q$

$$\tau_{v,w}(v, w) = v \otimes w$$

$$\tau'_{v,w}(v, w) = (-1)^{pq} v \otimes w \quad \leftarrow \text{Koszul's sign change.}$$

5) Δ^+ is not symmetric. (Idea: $\tau_{v,w}$ would have to be id.)
id \circ f \otimes g \neq g \otimes id !!!
↑ in graded

Def: (Symmetric Monoidal Functor)

For symmetric monoidal categories (V, \otimes, I, τ) and $(V', \otimes', I', \tau')$

a functor $F: V \rightarrow V'$ is a

symmetric monoidal functor if it is monoidal and for $X, Y \in \text{ob}(V)$

$$\tilde{\tau}_{X, Y} \circ F = \tilde{\tau}'_{F(X), F(Y)}$$

eg Extension of scalars (as before).

Def: (Commutative Monoid in Sym. Mon. Cat. V)

Let $(V, \otimes, I, \tilde{\tau})$ sym. mon. cat.

If $(M, \mu, \eta) \in \text{ob}(\underline{\text{Mon}}(V))$ and

$$\begin{array}{ccc} M \otimes M & \xrightarrow{\tilde{\tau}} & M \otimes M \\ \mu \searrow & & \swarrow \mu \\ & M & \end{array}$$

is commutative. Then we say M is a commutative monoid.

Defⁿ: (Category of Commutative Monoids
in \mathcal{V})

The class of commutative objects forms a category with morphisms given by monoid morphisms.

Denote this category by $cMon(\mathcal{V})$

Defⁿ: $\mathbb{F} =$ cat. of finite sets
with morphisms given
by fn.

Remark: $(\mathbb{F}, \sqcup, \emptyset, \tilde{c})$ is a symmetric
monoidal cat.

Theorem 2: $SymMonCat(\mathbb{F}, \mathcal{V}) \cong cMon(\mathcal{V})$

Aim: Define Frobenius objects and find the analogue of Δ^+ and Φ to make a thm 3.

Def: (Frobenius Objects)

A Frobenius Object in a monoidal category (V, \otimes, I) is an object $A \in \text{Ob}(V)$ with four maps

$$\eta: I \rightarrow A \quad \mu: A \square A \rightarrow A \quad \delta: A \rightarrow A \square A \quad \varepsilon: A \rightarrow I$$

$$s.t) \quad \mu \circ (\text{id}_A \square \eta) = \underset{\substack{\uparrow \\ \text{proj onto } A}}{\text{id}_A} = \mu (\eta \square \text{id}_A) \quad (\text{multiplication and unit})$$

$$2) \quad (\text{id}_A \square \varepsilon) \circ \delta = \text{id}_A = (\varepsilon \square \text{id}_A) \circ \delta \quad (\text{comulti and counit})$$

$$3) \quad (\text{id}_A \square \mu) \circ (\delta \square \text{id}_A) = \delta \circ \mu = (\mu \square \text{id}_A) \circ (\text{id}_A \square \delta) \quad (\text{Frobenius relation})$$

We denote the collection of Frobenius objects by $\text{Frob}(V)$.

Ex: In $\text{Vect}_{\mathbb{K}}$ the Frobenius objects are given by the Frobenius algebras.

See 2.3.24 of Kock

Def: Frobenius homomorphisms are morphisms between Frobenius objects preserving the structures $\mu, \eta, \delta, \varepsilon$.

Remark: $\text{Frob}(V)$ is a category with Frobenius morphisms.

Def: (Commutative Frobenius Objects)

A commutative Frobenius object in a symmetric monoidal category (V, \otimes, I, α) is commutative if

$$\mu \circ \tilde{\tau}_{x,x} = \mu$$

and cocommutative if

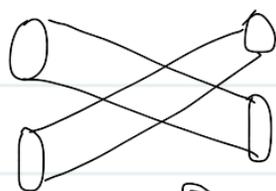
$$\tilde{\tau}_{x,x} \circ \delta = \delta$$

Lemma: A Frobenius object is commutative if it is cocommutative.

Remark: The set of commutative Frobenius objects $\text{cFrob}(V)$ forms a category with Frobenius morphisms.

Next: 2Cob.

2Cob symmetric monoidal category



twists



every object is Frobenius.